Inhomogeneous hypergraphs

Élie de Panafieu,

Algebra Institute, J. Kepler University, Linz, Autriche

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Tool:	Analytic combinatorics
Model:	Inhomogeneous hypergraphs
Application:	Constraint Satisfaction Problems (CSP)





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$$\exists x_1, \dots, x_n, \ C_1\left(x_{r(1,1)}, \dots, x_{r(1,a_1)}\right) \land \dots \land C_m\left(x_{r(m,1)}, \dots, x_{r(m,a_m)}\right)$$

Enumeration of satisfied instances.

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Generating function

$$A(z) = \sum_{a \in \mathcal{A}} \frac{z^{|a|}}{|a|!},$$

combinatorial description

asymptotics and limit laws

→ equation on the generating function,

P)

 \leftarrow analytic properties.

Inhomogeneous graphs

"The evolution of graphs may be considered as a rather simplified model of the evolution of certain communication nets [...]. Of course, if one aims at describing such a real situation, one should replace the hypothesis of equiprobability of all connections by some more realistic hypothesis."

– Erdős, Rényi (1959)

Probabilistic model introduced by Söderberg, extended by Bollobás, Janson, Riordan. Enumerative model analyzed by E.d.P. and Ravelomanana.

- The model inputs a parameter $R \in \operatorname{Sym}_q(\mathbb{R}_{\geq 0})$,
- each vertex v receives a color t(v) in $\{1, \ldots, q\}$,
- each edge (v, w) has a weight $R_{t(v),t(w)}$.

Inhomogeneous graphs are counted with a weight

weight(G) =
$$\prod_{(v,w) \in edges(G)} R_{t(v),t(w)}.$$

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Example: properly q-colored graphs

A graph is properly *q*-colored if each vertex has a color in $\{1, \ldots, q\}$, and no edge links two vertices having the same color. Bijection between inh. graphs and properly *q*-colored graphs

Asymptotics for n vertices and m edges

$$\frac{n^{2m}}{2^m m!} e^{\left(\frac{m}{n}\right)^2 \frac{q}{q-1}} \left(1 + \frac{2}{q-1} \frac{m}{n}\right)^{-\frac{q-1}{2}} \left(1 - \frac{1}{q}\right)^m q^n \left(1 + o(1)\right)$$

when m/n has a positive limit. A result already obtained by Wright (1972).

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$$R = \begin{pmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix} \qquad G = \underbrace{\begin{array}{c} & 1 \\ 0 & -1 \\ 1 &$$

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2-colorable graphs

A 2-colorable graph with c components has 2^c proper colorations.

Structure of inh. graph \rightarrow enumeration of 2-colorable graphs.



result already obtained by Pittel and Yeum (2004).

Example: friendship graphs

Each vertex has r hobbies among a set of size s, two vertices can be linked only if they share at least t hobby.



(r, s, t) = (3, 10, 1).

R is the adjacency matrix of the complement of the Kneser graph.

Example: systems of 2-equations

Consider a finite set *E* of triplets from \mathbb{F}_d , and a system

$$\exists x_1, \ldots, x_n \in \mathbb{F}_d, \begin{cases} a_1 x_{e(1,1)} + b_1 x_{e(1,2)} &= c_1, \\ & \vdots \\ a_m x_{e(m,1)} + b_m x_{e(m,2)} &= c_m. \end{cases}$$

The probability of satisfiability is expressed using the enumeration of inhomogeneous graphs



For 2QXorSAT, R is a Hamming-like matrix.

Constraint Satisfaction Problems

A CSP is a set of Boolean functions taking value in a finite set. An instance is a formula

 $\exists x_1, \ldots, x_n, \ C_1(x_{r(1,1)}, \ldots, x_{r(1,a_1)}) \land \cdots \land C_m(x_{r(m,1)}, \ldots, x_{r(m,a_m)})$

A satisfied instance is a pair instance-solution.

Bijection between satisfied instances of CSP with clauses of arity 2, and inhomogeneous graphs

variable	vertex,
constraint	edge,
values	colors,
number of clauses satisfied	$R_{i,i}$.
by (i,j) or (j,i)	

Inhomogeneous hypergraphs

- Each vertex v receives a color t(v) in $\{1, \ldots, q\}$,
- edges can contain more than 2 vertices,
- the weight of an edge depends of the colors it contains.



 $\mathsf{orderings}(\mathit{G}) = |\{((5,2,3),(4,1),(2,3),(1,5,3),(1,6)),\ldots\}|,$

weight(G) =
$$\frac{\text{orderings}(G)}{|E(G)|! \prod_{e \in E(G)} |e|!} \prod_{e \in E(G)} \omega_{\overline{t}(e)}$$

Many parameters: for all $\overline{t} \in \mathbb{N}^q$, weight $\omega_{\overline{t}}$ for the edges that contain t_i vertices of type i.

$$\Omega(x_1,...,x_q) = \sum_{t_1,...,t_q \ge 0} \omega_{t_1,...,t_q} \frac{x_1^{t_1}}{t_1!} \cdots \frac{x_q^{t_q}}{t_q!}.$$



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$$T_i(z) = z e^{\partial_i \Omega(T_1(z),...,T_q(z))}$$



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Trees and unicycles

Unrooted trees: $U(z) = \sum_{i} T_{i} + \Omega(\bar{T}) - \bar{T}\bar{\partial}\Omega(\bar{T})$ Unicycles: $V(z) = -\frac{1}{2}\log\left(\det\left(\operatorname{Id} - \operatorname{diag}(\bar{T})\mathcal{H}_{\Omega}(\bar{T})\right)\right)$ Trees and unicycles: excess $k = \sum_{e \in \operatorname{edges}}(|e| - 1) - n$

$$n![z^n]\frac{U(z)^{-k}}{(-k)!}e^{V(z)} = \frac{n!}{2i\pi} \oint \frac{U(z)^{-k}}{(-k)!}e^{V(z)}\frac{dz}{z^{n+1}}$$



(plot from Flajolet, Sedgewick 2009)

Sum of the weights of inhygraphs with n vertices and excess k

$$[y^{k+n}] \sum_{n_1+\dots+n_q=n} \binom{n}{n} \prod_{\overline{t}\in\mathbb{N}^q} (1+\omega_{\overline{t}}y^{t_1+\dots+t_q-1})^{\binom{n_1}{t_1}\dots\binom{n_q}{t_q}}$$

Approximations: coefficient extraction, Stirling, integral for the sum

$$\sim C_{n,k} \int_{\substack{x_1+\cdots+x_q=1\\ar{x}\in[0,1]^q}} A(ar{x}) e^{n\Phi(ar{x})} dar{x}$$

where $\Psi(\bar{x}) = \frac{\bar{x}\bar{\partial}\Omega(\bar{x}) - \Omega(\bar{x})}{\sum_{i=1}^{q} x_i}$, $\Psi(\zeta \bar{x}) = 1 + \frac{k}{n}$ and $\Phi(\bar{x}) = \bar{x} \left(\log(\bar{\nu}) - \log(\bar{x})\right) + \frac{\Omega(\zeta \bar{x})}{\zeta} - \left(1 + \frac{k}{n}\right)\log(\zeta)$.



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Forthcoming research

Applications will provide new inhygraphs properties to investigate

- structure of graphs with a giant component generating function of connected graphs with the same density of edges,
- graphs with forbidden subgraphs satisfiability threshold of 2-SAT,
- data base modelling using inhomogeneous hypergraphs analysis of data mining algorithms (Peps HYDrATA).